

# Quantum Computation of Prime Number Functions

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Joint workshop:

Quantum Physics: from fundamental questions to applications  
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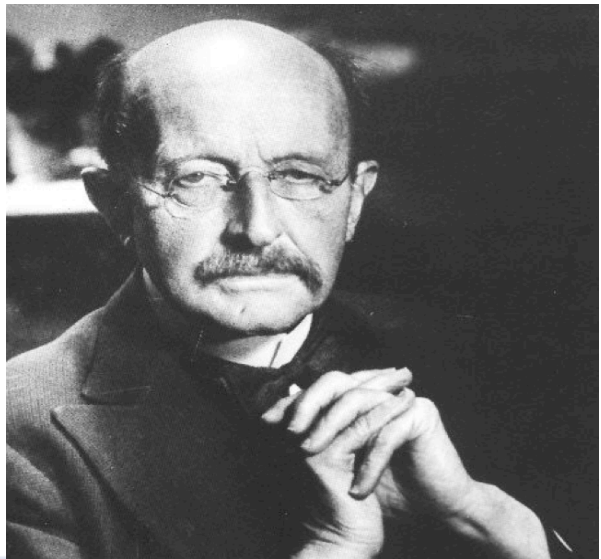
# CLASSICAL PHYSICS



# MATHEMATICS

IR

# QUANTUM PHYSICS

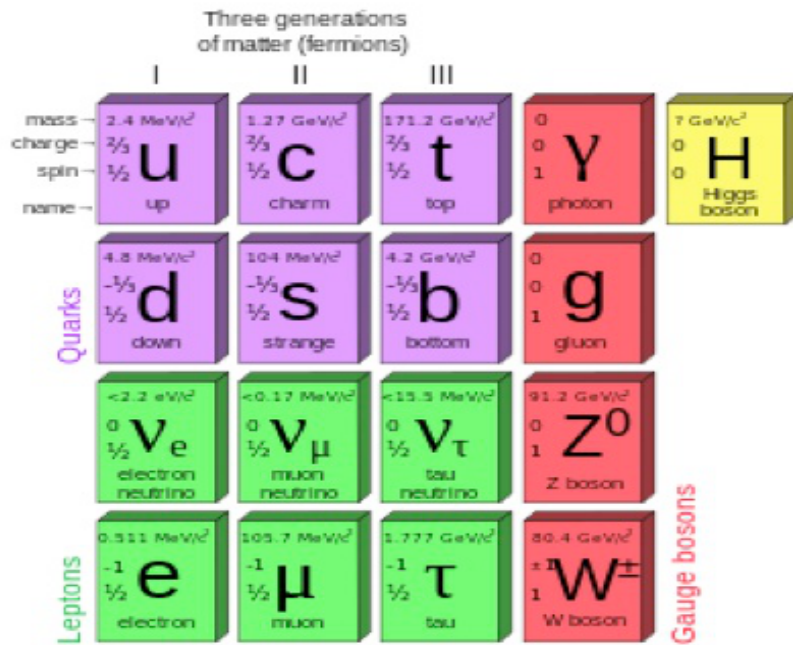


N

# Fundamental building blocks

NATURE

NUMBERS



1	2	3	4	5	6	7	8	9	10
11	12	13	14	15	16	17	18	19	20
21	22	23	24	25	26	27	28	29	30
31	32	33	34	35	36	37	38	39	40
41	42	43	44	45	46	47	48	49	50
51	52	53	54	55	56	57	58	59	60
61	62	63	64	65	66	67	68	69	70
71	72	73	74	75	76	77	78	79	80
81	82	83	84	85	86	87	88	89	90
91	92	93	94	95	96	97	98	99	100

In some sense quantum theory is a bending of physics towards number theory. However, deep facts of number theory play no role in questions of quantum mechanics....

In particular we do not know of any fundamental physical theories that are based on deep facts in number theory.

I would think that quantum mechanics will be completely reformulated and that number theory will play a key role in this formulation.

C. Vafa (2000)

While we wait for this reformulation  
let us see if Quantum Mechanics  
can do something for Number Theory

## Classical computer

n bits  $x = x_0 2^0 + x_1 2^1 + \dots + x_{n-1} 2^{n-1}$ ,  $x_i = 0, 1$ ,  $x = 0, 1, \dots, 2^n - 1$

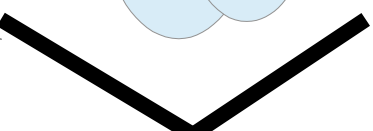
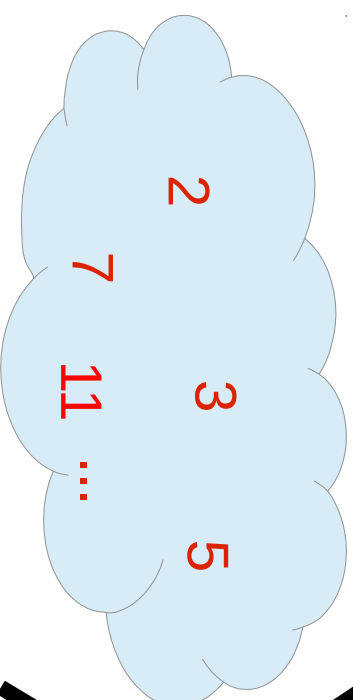
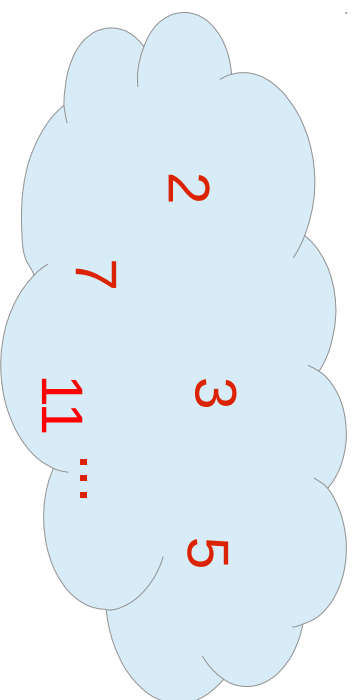
## Quantum computer

n qubits  $|x\rangle = |x_{n-1}, \dots, x_0\rangle = |x_{n-1}\rangle \otimes \dots \otimes |x_0\rangle$

Primes



State



## The Prime State

$$|P(n)\rangle = \frac{1}{\sqrt{\pi(2^n)}} \sum_{p < 2^n \in \text{Primes}} |p\rangle$$

$\pi(2^n)$  is the prime counting function

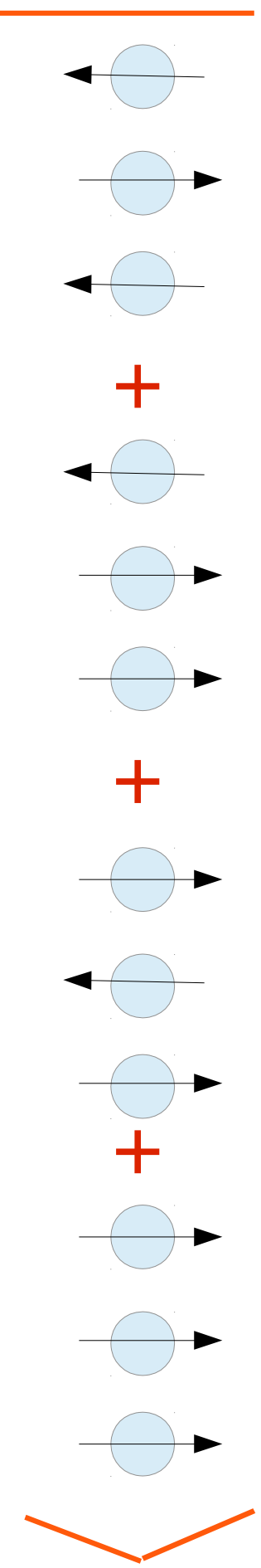


Quantum Mechanics allows for the superposition of primes implemented as states of a computational basis

$$|P(n)\rangle = \frac{1}{\sqrt{\pi(2^n)}} \sum_{p < 2^n \in \text{Primes}} |p\rangle$$

EX.  $n=3$

$$|P(3)\rangle = \frac{1}{\sqrt{4}} (|2\rangle + |3\rangle + |5\rangle + |7\rangle)$$



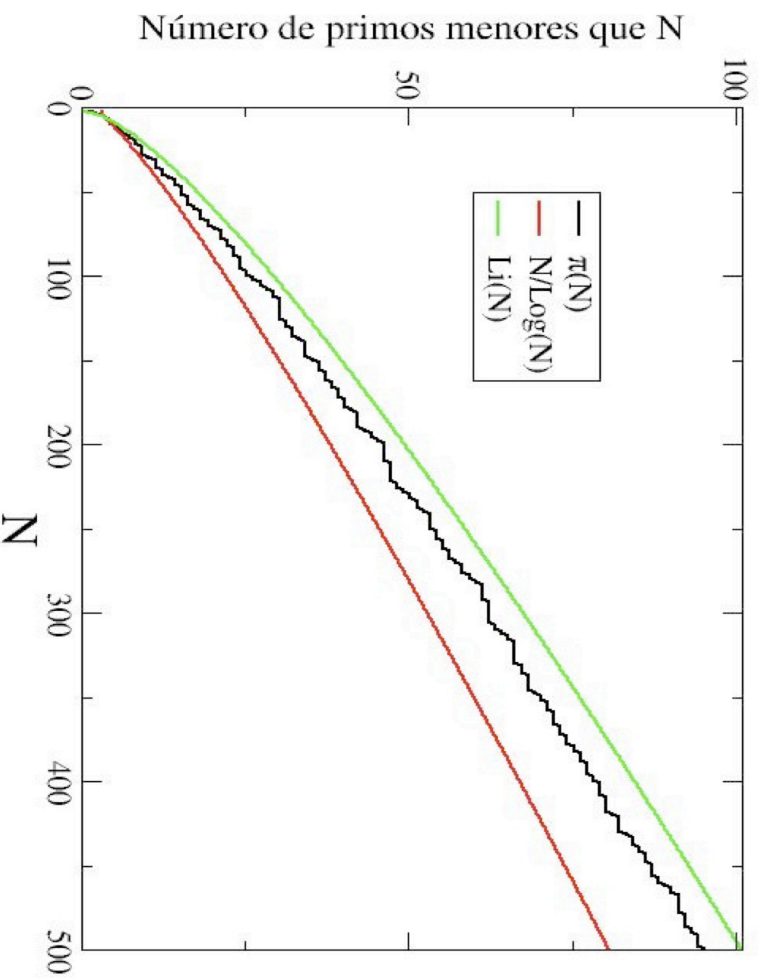
## Prime counting function

$$\pi(x)$$

Number of primes  $p$  less than or equal to  $x$  e.g.

$$\pi(100) = 25$$

Gauss – Legendre law



$$\pi(x) \approx \text{Li}(x) \approx \frac{x}{\ln x}$$

$$x \rightarrow \infty$$

Prime number theorem

- Hadamard (1896)
- de la Vallée-Poussin

## Prime Number Theorem (PNT)

$$\pi(x) \approx Li(x) \quad Li(x) = \int_2^x \frac{dt}{\log t} \approx \frac{x}{\log x} + \frac{x}{\log^2 x} + \dots$$

Density of primes:

$$\frac{d\pi(x)}{dx} \approx \frac{1}{\ln x}$$

Largest known value  $\pi(10^{24}) = 18\,435\,599\,767\,349\,200\,867\,886 \approx 1.8 \cdot 10^{22}$

Platt (2012)

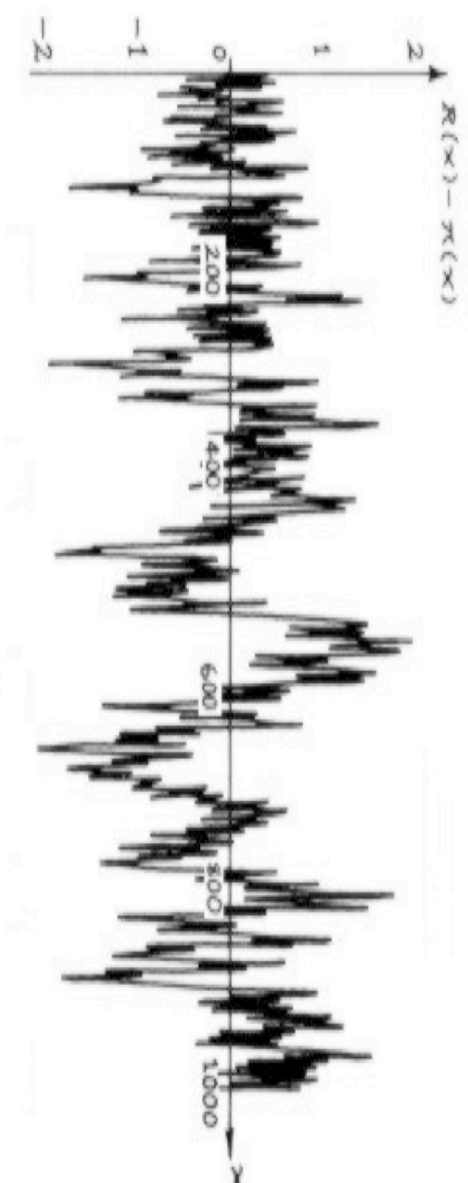
$$Li(10^{24}) - \pi(10^{24}) \approx 1.7 \cdot 10^{10}$$

The prime number function will oscillate around the Log Integral infinitely many times  
Littlewood, Skewes

A first change of sign is expected for some  $X < e^{727.9513468} \dots$

If the **Riemann hypothesis (RH)** is correct, fluctuations are bounded

$$|Li(x) - \pi(x)| < \frac{1}{8\pi} \sqrt{x} \log x$$



$$\zeta(s) = \sum_{n \geq 1} \frac{1}{n^s}$$

If  $\zeta(s) = 0$  with  $0 \leq \text{Real}(s) \leq 1$  then  $\text{Real}(s) = \frac{1}{2}$

Could the Prime state be constructed?

Does it encode properties of prime numbers?

What are its entanglement properties?

Could it provide the means to explore Arithmetics?

# Entanglement: single qubit reduced density matrices

$$|P(n)\rangle = \frac{1}{\sqrt{\pi(2^n)}} \sum_{i_{n-1}, \dots, i_1, i_0=0,1} P_{i_{n-1}, \dots, i_1, i_0} |i_{n-1}, \dots, i_1, i_0\rangle$$

$$P_{i_{n-1}, \dots, i_1, i_0} = \begin{cases} 1 & p = i_{n-1} 2^{n-1} + \dots + i_0 = \text{prime} \\ 0 & \text{otherwise} \end{cases}$$

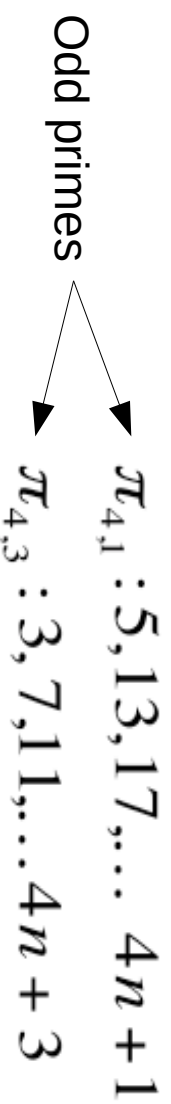
Density matrix qubit  $i=1$

$$\rho_{ab}^{(1)} = \frac{1}{\pi(2^n)} \sum_{i_{n-1}, \dots, i_2, i_0=0,1} P_{i_{n-1}, \dots, i_2, a, i_0} P_{i_{n-1}, \dots, i_2, b, i_0}$$

$$\rho_{00}^{(1)} = \frac{\pi_{4,1}(2^n)}{\pi(2^n)}$$

$$\rho_{11}^{(1)} = \frac{1 + \pi_{4,3}(2^n)}{\pi(2^n)}$$

$$\rho_{01}^{(1)} = \frac{\pi_2^{(1)}(2^n)}{\pi(2^n)}$$



### Dirichlet theorem:

There infinite number of primes of the form  $1 + 4n$  and  $3 + 4n$

### PNT for arithmetic series

$$\lim_{x \rightarrow \infty} \frac{\pi_{4,1}(x)}{Li(x)} = \lim_{x \rightarrow \infty} \frac{\pi_{4,3}(x)}{Li(x)} = \frac{1}{\phi(4)} = \frac{1}{2} \longrightarrow S(\rho^{(i)}) \sim \log 2$$

### Chebyshev bias:

For low values of  $x$  there exist more primes  $1 \pmod{4}$  than  $3 \pmod{4}$

$$\Delta(x) = \pi_{4,3}(x) - \pi_{4,1}(x)$$

Related to magnetization of qubit  $i=1$

$$\langle \sigma_z^{(1)} \rangle = \frac{-\Delta(2^n) - 1}{\pi(2^n)}$$

Twin primes :  $p, p+2$

$$\text{Counting function} \quad \pi_2(x) \approx 2C_2 \frac{x}{(\log x)^2} \quad (\text{Hardy-Littlewood conjecture})$$

Twin primes  $\swarrow \searrow$   
 $\pi_2^{(1)} : (5, 7), \dots (1 \pmod 4, 3 \pmod 4)$   
 $\pi_2^{(3)} : (11, 13), \dots (3 \pmod 4, 1 \pmod 4)$

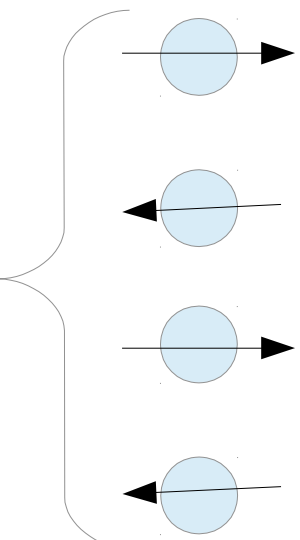
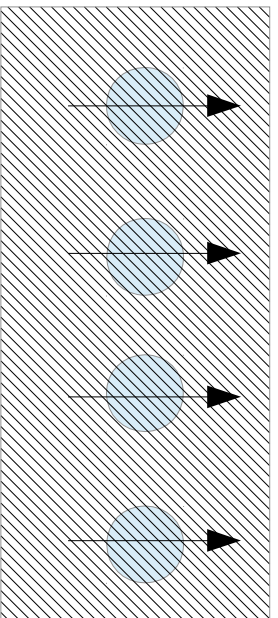
$$\langle \sigma_x^{(1)} \rangle = \frac{2\pi_2^{(1)}(2^n)}{\pi(2^n)}, \quad \langle \sigma_x^{(1)} \sigma_x^{(2)} + \sigma_y^{(1)} \sigma_y^{(2)} \rangle = \frac{4\pi_2^{(3)}(2^n)}{\pi(2^n)}$$

Twinsip  $\rightarrow$  off diagonal entries of density matrix

Sub-series of primes, twin primes, etc. are amenable to measurements

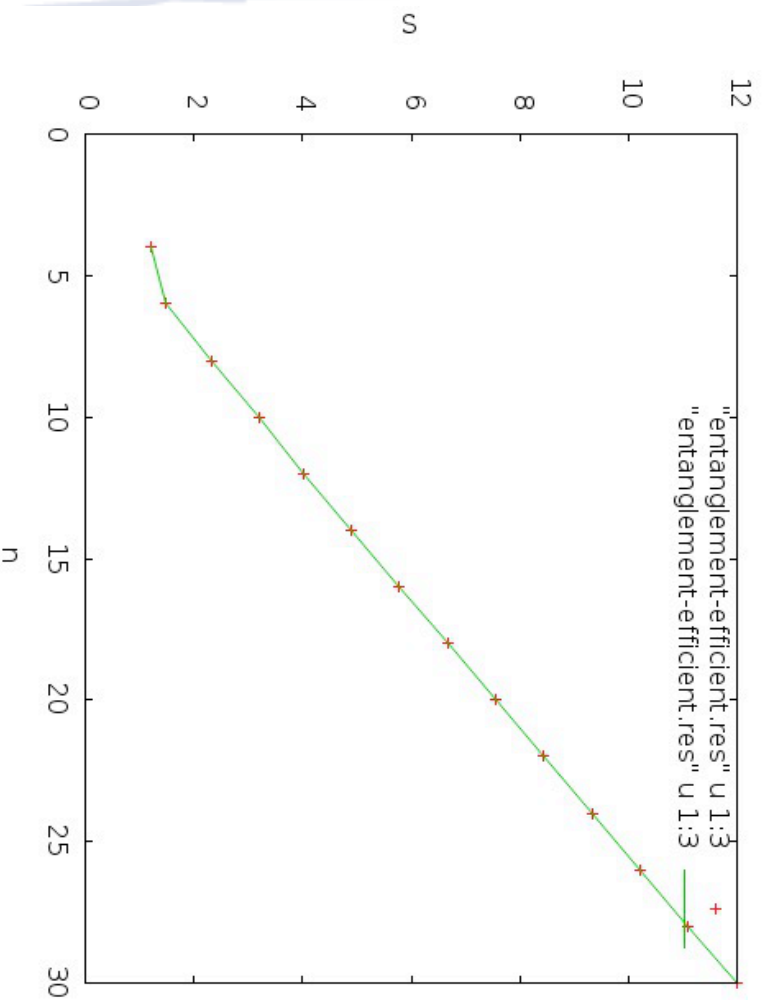


# Entanglement entropy of the Prime state



$$\rho_n$$

$$\frac{n}{2}$$



“There is entanglement in the Primes”

Volume law scaling

$$S \sim .8858 n + \text{const}$$

A. Monras, G. Sierra, JIL

## Scaling of entanglement entropy

$$S \sim n - \text{const}$$

Random states

$$S \sim .8858 n + \text{const}$$

Prime state

$$S \sim n^{\frac{d-1}{d}} + \text{const}$$

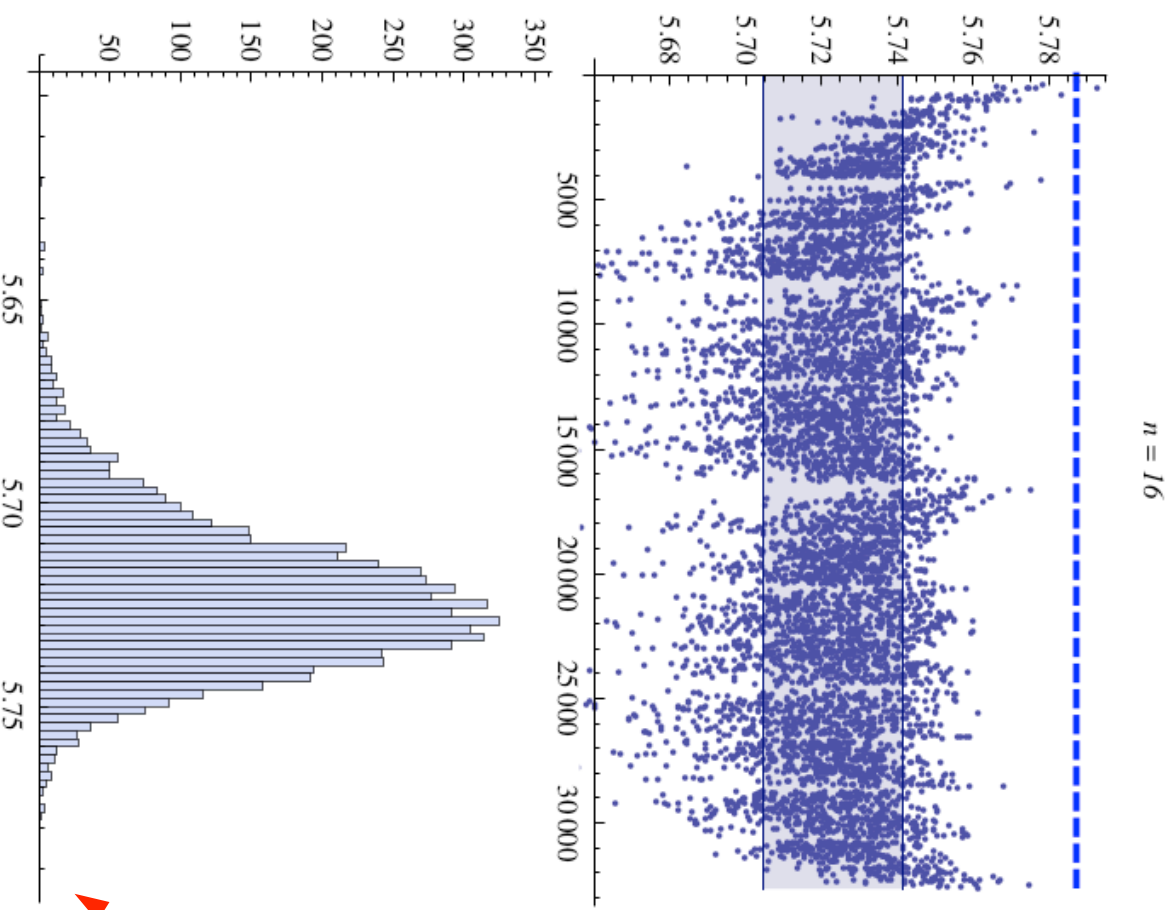
Area law in d-dimensions

$$S \sim \frac{c}{3} \log n + \text{const}$$

Critical scaling in d=1  
at quantum phase transitions

$$S \sim \log(\xi) = \text{const}$$

Finitely correlated states  
away from criticality

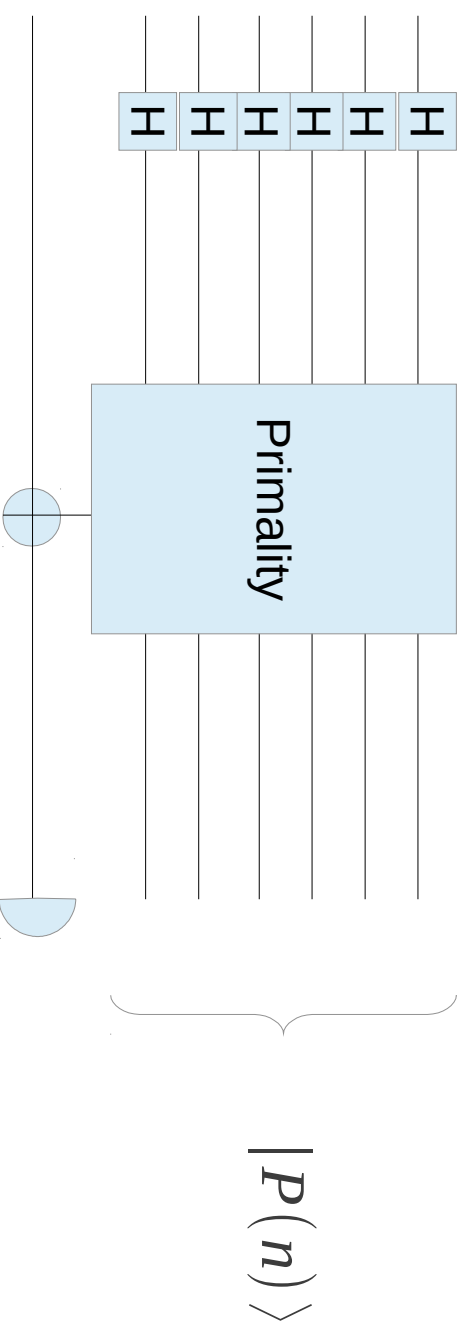


Entanglement entropy  
for different partitions

Original partition carries  
more entanglement (!?)

$$\mu = 5.72307; \quad \sigma = 0.0184293; \quad \frac{S_{\max} - \mu}{\sigma} = 3.47743$$

# Construction of the Prime state



$$U_{\text{primality}} \sum_x |x\rangle |0\rangle = |P(n)\rangle |0\rangle + \sum_{c \in \text{composite}} |c\rangle |1\rangle$$

$$\text{Prob}(|P(n)\rangle) = \frac{\pi(2^n)}{2^n} \approx \frac{1}{n \log 2}$$

Efficient construction

## Construction of twin primes

$$U_{+2} |P(n)\rangle |0\rangle = \sum_{p \in \text{primes}} |p+2\rangle |0\rangle$$

$$U_{\text{primality}} U_{+2} |P(n)\rangle |0\rangle = \sum_{p, p+2 \in \text{primes}} |p+2\rangle |0\rangle + \sum_{p+2 \notin \text{primes}} |p+2\rangle |1\rangle$$

$$\Pr(\text{twin primes}) = \frac{\pi_2(2^n)}{\pi(2^n)} \approx \frac{2C_2}{n \log 2}$$

# Grover construction of the Prime state

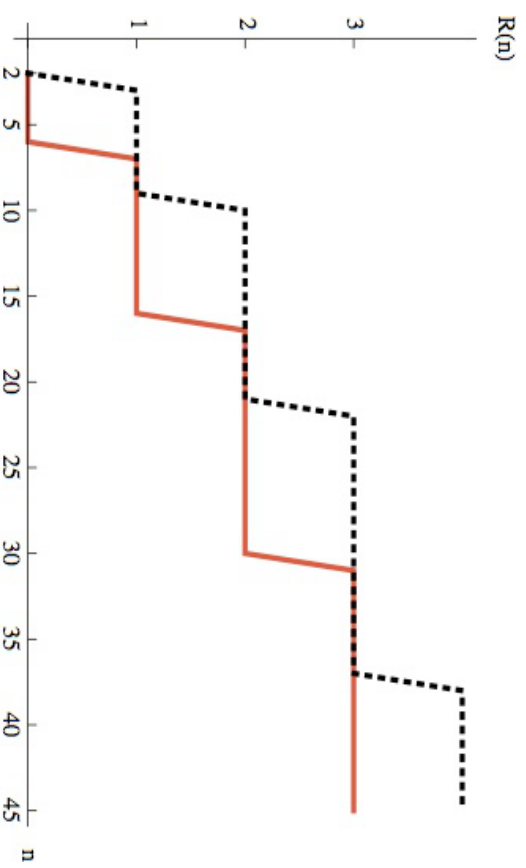
$$|\psi_0\rangle = \sum_{x < 2^n} |x\rangle = \frac{1}{\pi(2^n)} \left( \underbrace{\sum_{p \in \text{primes}} |p\rangle}_{M} + \underbrace{\sum_{c \in \text{composites}} |c\rangle}_{N} \right)$$

# calls to Grover

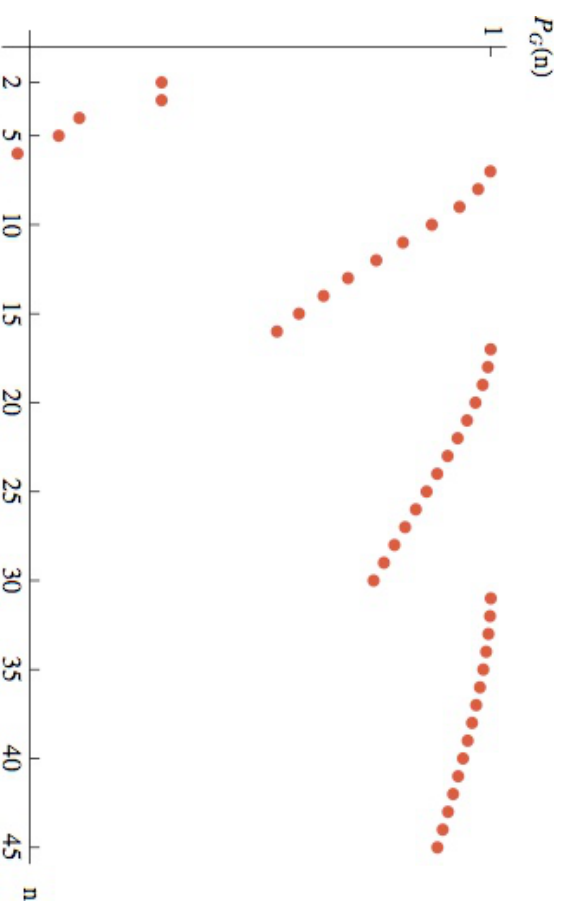
$$R(n) \leq \frac{\pi}{4} \sqrt{\frac{N}{M}} \leq \frac{\pi}{4} \sqrt{n \log 2}$$

$$|\psi_f\rangle = |P(n)\rangle$$

# calls to Grover



Overlap between  
Grover state and the Prime state



We need to construct an oracle!

# Construction of a Quantum Primality oracle

An efficient Quantum Oracle can be constructed using classical primality tests

## Miller-Rabin primality test

- Find  $s$  and  $d$  (odd) such that
$$x - 1 = 2^s d$$
- Choose witness  $a$ 
$$1 \leq a \leq x$$
- If  $a^d \not\equiv 1 \pmod{x}$  then  $x$  is composite with certainty
$$a^{2^r d} \not\equiv -1 \pmod{x} \quad 0 \leq r \leq s-1$$
- If the test fails,  $x$  may be prime or composite.
- Latter case:  $a$  is a strong liar to  $x$
- Eliminate strong liars checking less than  $\log^2 x$  witnesses



# Construction of a Quantum Primality oracle


An efficient Quantum Oracle can be constructed using classical primality tests

## Miller-Rabin primality test

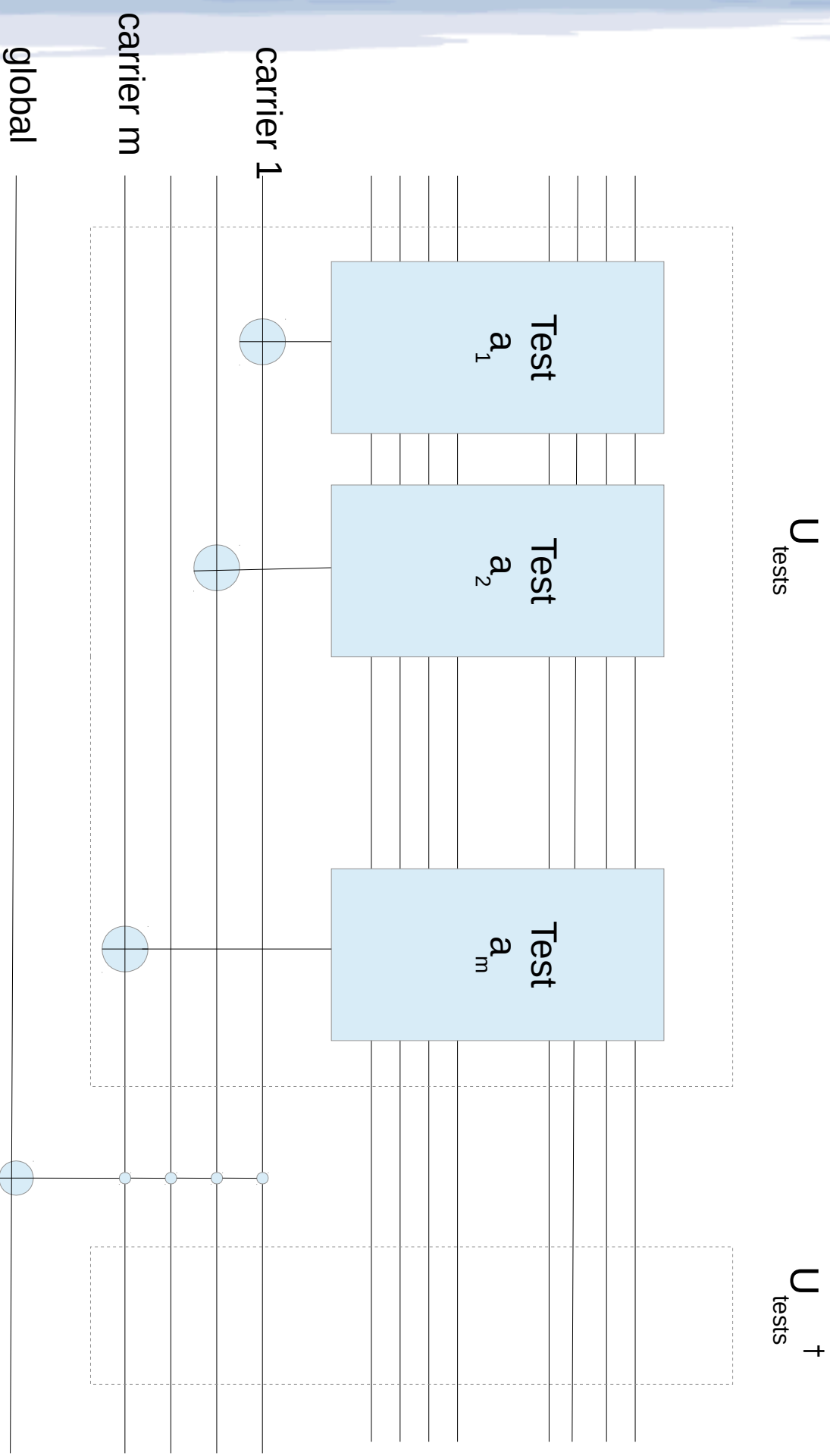
- Find  $s$  and  $d$  (odd) such that
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Finding d and s

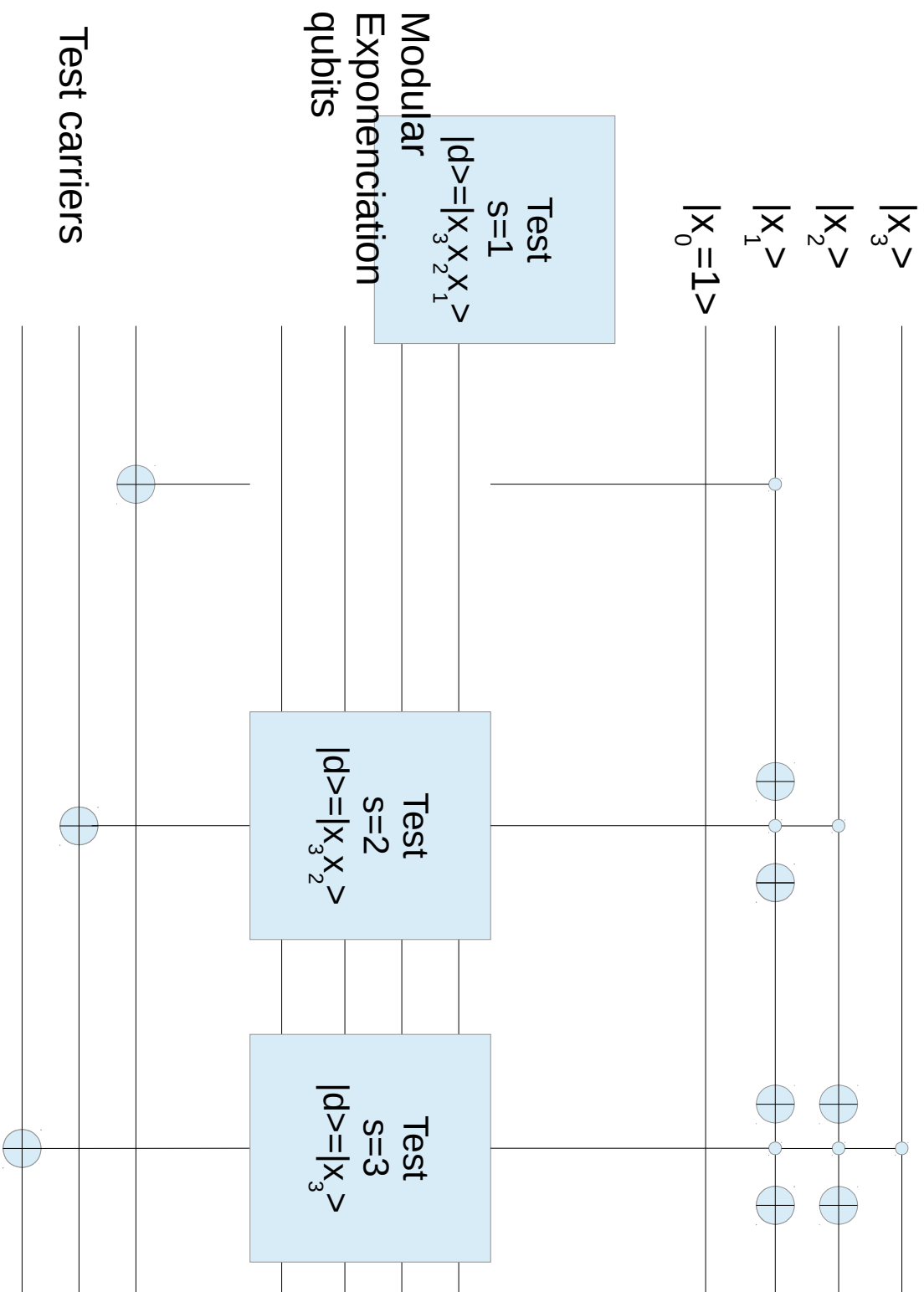
$$x = 49 \rightarrow x - 1 = 48 = 2^4 \times 3 \rightarrow s = 4, d = 3$$

$$|49\rangle = |1,1,0,0,0,1\rangle \rightarrow |1,1,0,0,0,0\rangle$$


d s



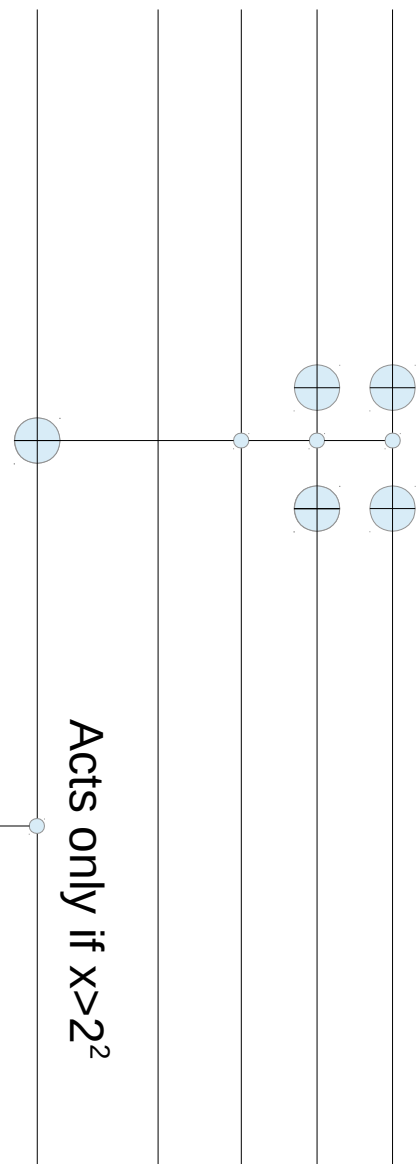
Structure of the quantum primality oracle



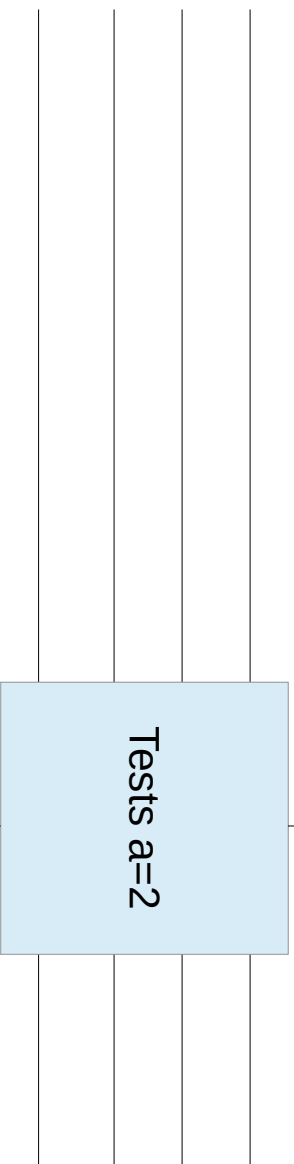
Tests are condition to the actual value of x

Is  $x < 2^2$  ?

$|x_3\rangle$   
 $|x_2\rangle$   
 $|x_1\rangle$   
 $|x_0\rangle$   
 $|1\rangle$



Modular  
Exponentiation  
qubits



Test carrier

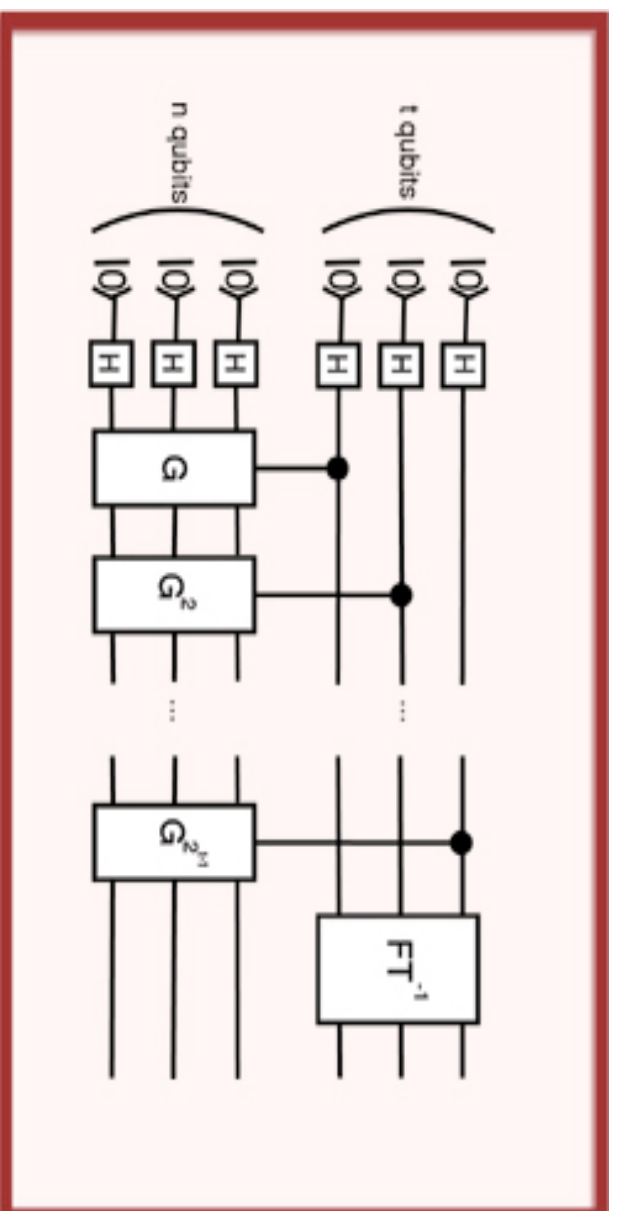


Tests are only run when witness is smaller than  $x$

# Quantum Counting of Prime numbers

quantum primality oracle + quantum counting algorithm

Brassard, Hoyer, Tapp (1998)



Counts the number of solutions to the oracle

We want to count  $M$  solutions out of  $N$  possible states

We know an estimate  $\tilde{M}$

$$|\tilde{M} - M| < \frac{2\pi}{c} M^{\frac{1}{2}} + \frac{\pi^2}{c^2}$$

Bounded error in quantum counting

Bounded error in the quantum counting of primes

$$\left| \pi_{QM}(x) - \pi(x) \right| \leq \frac{2\pi}{c} \frac{x^{1/2}}{\log^{1/2} x}$$

We use the  
PNT

$$\left| \pi_{QM}(x) - \pi(x) \right| \leq \frac{2\pi}{c} \frac{x^{1/2}}{\log^{1/2} x}$$

Error of counting is smaller than the bound for the fluctuations if Riemann hypothesis is correct

$$\frac{2\pi}{c} \frac{x^{\frac{1}{2}}}{\log^{\frac{1}{2}} x} < x^{\frac{1}{2}} \log x$$

Best classical algorithm by Lagarias-Miller-Odlyzko (1987) implemented by Platt (2012)

$$T \sim x^{\frac{1}{2}}$$

$$S \sim x^{\frac{1}{4}}$$

**A Quantum Computer could calculate the size of fluctuations more efficiently than a classical computer**

$$T \propto x^{\frac{1}{2}}$$

$$S \propto \log x$$



## Beyond Prime numbers: the **q**-functor

$$S \subseteq X \rightarrow |S\rangle = \frac{1}{\sqrt{|S|}} \sum_{x \in S} |x\rangle$$

$$|S\rangle = \frac{1}{\sqrt{|S|}} \sum_{x \in X} \chi_S(x) |x\rangle \quad \chi_S(x) = \begin{cases} 1 & x \in S \\ 0 & x \notin S \end{cases}$$

Primes

Average (Cramér) primes

Dirichlet characters

...

Only needs the construction of a quantum oracle for  $\chi_S(x)$

# Conclusion

Quantum Simulation of Arithmetics

Superposition of series of numbers using appropriate q-oracles

Measurements of arithmetic functions

More efficient approaches are likely

**Thank you**

**Gracies**

What is this?



**Gaudi Magic Square  
in Sagrada Familia  
Barcelona**